

## ALTERNATING HAMILTON CYCLES WITH MINIMUM NUMBER OF CROSSINGS IN THE PLANE

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### ABSTRACT

Let  $X$  and  $Y$  be two disjoint sets of points in the plane such that  $|X| = |Y|$  and no three points of  $X \cup Y$  are on the same line. Then we can draw an alternating Hamilton cycle on  $X \cup Y$  in the plane which passes through alternately points of  $X$  and those of  $Y$ , whose edges are straight-line segments, and which contains at most  $|X| - 1$  crossings. Our proof gives an  $O(n^2 \log n)$  time algorithm for finding such an alternating Hamilton cycle, where  $n = |X|$ . Moreover we show that the above upper bound  $|X| - 1$  on crossing number is best possible for some configurations. <sup>a</sup>.

*Keywords:* two point sets, plane, crossing number, Hamilton cycle, line embedding.

### 1. Introduction

Let  $X$  and  $Y$  be two disjoint sets of points in the plane such that no three points of  $X \cup Y$  are on the same line. Then there are some results on drawing graphs on  $X \cup Y$  whose edges are straight-line segments in the plane. For example, Tokunaga [3] showed that the minimum number of crossings of two spanning trees on  $X$  and on  $Y$  is determined by the number of edges of  $\text{conv}(X \cup Y)$  joining points of  $X$  to those of  $Y$ . Abellanas et al. [1] considered some problems on alternating spanning trees on  $X \cup Y$  with no crossings, whose edges join points of  $X$  to those of  $Y$ , and proved that there exists an alternating spanning tree on  $X \cup Y$  with no crossings whose maximum degree is at most  $O(|Y|/|X| + \log |X|)$ , where  $|Y| \geq |X|$ . Kaneko has recently showed that if  $|Y| = |X|$ , then there exists an alternating spanning tree

on  $X \cup Y$  with no crossings whose maximum degree is at most three. Moreover, Akiyama and Urrutia [2] gave an algorithm that determines if  $X \cup Y$  admits an alternating Hamilton path with no crossings under the assumption that  $|X| = |Y|$  and  $X \cup Y$  are the set of vertices of  $\text{conv}(X \cup Y)$ .

In this paper, we deal with an *alternating Hamilton cycle* on  $X \cup Y$ , which passes through all the points of  $X \cup Y$  such that points of  $X$  and  $Y$  are alternating and whose edges are straight-line segments. Our aim is to draw an alternating Hamilton cycle on  $X \cup Y$  in the plane so that it has a small number of crossings. In particular, we have the following question: What is the minimum number of crossings in such alternating Hamilton cycles? An answer to this question is given in the following theorem.

**Theorem 1** *Let  $X$  and  $Y$  be two disjoint sets of points in the plane such that  $|X| = |Y|$  and no three points of  $X \cup Y$  are on the same line. Then we can draw an alternating Hamilton cycle on  $X \cup Y$  which has at most  $|X| - 1$  crossings (see Figure 1). Moreover there exist configurations  $X \cup Y$  for which this upper bound  $|X| - 1$  is best possible.*

Since our proof of the above theorem is constructive, it provides an algorithm for drawing an alternating Hamilton cycle given in the theorem. Actually, by a direct analysis of the proof, we can draw it in  $O(n^2 \log n)$  time, where  $n = |X|$ . Moreover, by the proof we can draw it so that no three edges of it intersect at the same point.

We first show that the upper bound  $|X| - 1$  on crossing number is best possible for some configurations. Let  $P$  be a convex polygon with  $2n$  vertices such that no three line segments joining two points of  $P$  intersect at the same point except their common end point, and let  $X = \{x_1, x_2, \dots, x_n\}$  and  $Y = \{y_n, y_{n-1}, \dots, y_1\}$  be disjoint sets of vertices of  $P$  such that  $x_1, x_2, \dots, x_n, y_n, y_{n-1}, \dots, y_1$  are consecutive vertices of  $P$ , in particular,  $X$  and  $Y$  can be separated by a certain straight-line in the plane (see Figure 1). Then we now show that every alternating Hamilton cycle on  $X \cup Y$  has at least  $n - 1$  crossings. Let  $C$  be any alternating Hamilton cycle on  $X \cup Y$ . Then  $C$  contains at least  $2n - 2$  edges  $x_i y_j$  which is neither  $x_1 y_1$  nor  $x_n y_n$ . Let  $x_i y_j$  be one of such edges. Since the path  $C - x_i y_j$  connects  $x_i$  and  $y_j$ , and  $P$  is a convex polygon, the edge  $x_i y_j$  must intersect at least one edge of  $C - x_i y_j$ , say  $x_r y_s$ , and we get a pair  $(x_i y_j, x_s y_t)$  of intersecting edges. Because there exist at least  $n - 1$  distinct such pairs of edges,  $C$  contains at least  $n - 1$  crossings. It is obvious that there exists an alternating Hamilton cycle on  $X \cup Y$  with  $n - 1$  crossings (see Figure 1). Therefore the upper bound in Theorem 1 is best possible for the above configurations.

## 2. Proof of Theorem 1

Let  $S$  be a set of points in the plane. Then we denote by  $\text{conv}(S)$  the *convex hull* of  $S$ , which is the smallest convex set containing  $S$ . We briefly say that  $S$  is *in general position* if no three points of  $S$  lie on the same line. The following lemma plays an important role.

**Lemma 1** *Let  $R$  and  $B$  be disjoint sets of points in the plane and  $x$  be a point in the plane not contained in  $\text{conv}(R \cup B)$  such that  $2 \leq |R| \leq |B|$  and  $R \cup B \cup \{x\}$  is in*

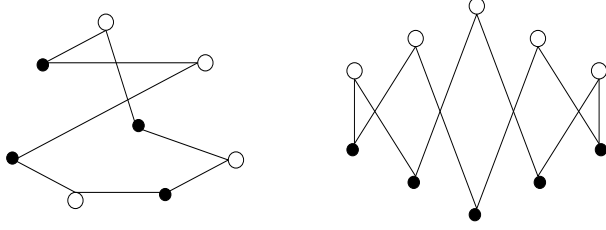


Fig. 1. Alternating Hamiltonian cycles with two and four crossings.

*general position. If the two vertices of  $\text{conv}(R \cup B \cup \{x\})$  adjacent to  $x$  are contained in  $R$ , then we can partition  $R \cup B$  into  $D_1 \cup D_2$  which satisfies the following three conditions: (i)  $|D_1 \cap R| = |D_1 \cap B|$ , (ii)  $\text{conv}(D_1 \cup \{x\})$  and  $\text{conv}(D_2 \cup \{x\})$  intersect only at  $x$ , and (iii) for each  $i \in \{1, 2\}$ , one vertex of  $\text{conv}(D_i \cup \{x\})$  adjacent to  $x$  is contained in  $B$  (see Figure 2).*

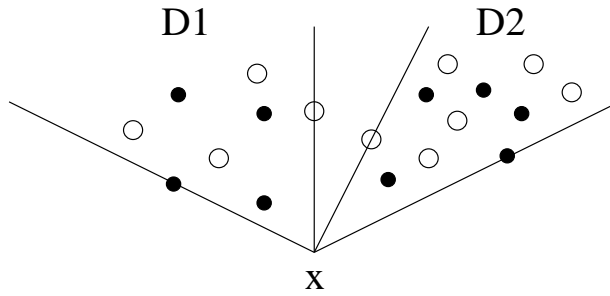


Fig. 2. A partition  $D_1 \cup D_2$  of  $R \cup B$ , where  $R = \{\bullet\}$  and  $B = \{\circ\}$ .

*Proof* We prove the lemma by induction on  $|R|$ . If  $|R| = 2$  then the lemma follows immediately, and so we may assume  $|R| \geq 3$ . By a suitable rotation of the plane, we may assume that  $x$  lies on the bottom of  $\text{conv}(R \cup B \cup \{x\})$ . Moreover in this case, we shall prove the lemma with the additional property that  $D_1$  lies to the left of  $D_2$ . We now consider rays emanating from  $x$  and going upward, and so a ray means such a ray in this proof. We define a function  $f$  of a ray  $r$  by

$$f(r) = \{\text{the number of points of } R \text{ lying on or to the left of } r\} \\ - \{\text{the number of points of } B \text{ lying on or to the left of } r\}.$$

Let  $r_1$  and  $r_2$  denote the rays that pass through the left vertex and the right vertex of  $\text{conv}(R \cup B \cup \{x\})$  adjacent to  $x$ , respectively, and let  $r'_2$  be a ray obtained from  $r_2$  by a very small counterclockwise rotation around  $x$ . Then  $f(r_1) = 1$  and  $f(r'_2) = |R| - 1 - |B| \leq -1$ . Since the value of  $f$  changes  $\pm 1$  and since  $f(r_1) > 0 > f(r'_2)$ , there exists a ray  $r_a$  such that  $f(r_a) = 0$  and  $r_a$  passes through a point of  $B$ . We next rotate  $r_a$  clockwise around  $x$  until it meets a new point of  $R \cup B$ , and denote the ray by  $r_b$ . If  $r_b$  passes through a point of  $B$ , then let  $D_1$  and  $D_2$  be the set of points of  $R \cup B$  lying on or to the left of  $r_a$  and that of those points lying on or to the right of  $r_b$ , respectively. Then  $D_1 \cup D_2$  is the desired partition.

If  $r_b$  passes through a point of  $R$ , then let  $R_1$  and  $B_1$  be the subsets of  $R$  and

$B$ , respectively, whose points lie on or to the right of  $r_b$ . Then  $R_1$ ,  $B_1$  and  $x$  satisfy the assumption of the lemma, and hence by the inductive hypothesis, there exists a partition  $D_3 \cup D_4$  of  $R_1 \cup B_1$  that satisfies the conditions of this lemma together with the condition that  $D_3$  lies to the left of  $D_4$ . Consequently we can obtain the desired partition  $(D_1 \cup D_3) \cup D_4$  of  $R \cup B$ .  $\square$

For two disjoint sets  $S$  and  $T$  of points in the plane, we can define an *alternating Hamilton path* on  $S \cup T$  in the same way as an alternating Hamilton cycle on  $S \cup T$ . Of course, if there exists an alternating Hamilton path on  $S \cup T$  which begins with a vertex in  $S$ , then  $|T| \leq |S| \leq |T| + 1$ . An alternating Hamilton path connecting two points  $x$  and  $y$  is called an alternating Hamilton  $x - y$  path and denoted by  $P(x, y)$ . In order to prove Theorem 1, we need the following lemma, which will be proved latter.

**Lemma 2** *Let  $R$  and  $B$  be disjoint sets of points in the plane such that  $R \cup B$  is in general position. Suppose that a point, say  $r$ , of  $R$  is a vertex of  $\text{conv}(R \cup B)$ . Then the following two statements hold.*

(i) *If  $|R| = |B|$ , then there exists an alternating Hamilton path on  $R \cup B$  which begins with  $r$  and ends with a vertex of  $B$ , and which has at most  $|R| - 1$  crossings. Furthermore, if a point, say  $b$ , of  $B$  is a vertex of  $\text{conv}(R \cup B)$ , then there exists an alternating Hamilton  $r - b$  path which has at most  $|R| - 1$  crossings.*

(ii) *If  $|R| = |B| + 1$ , then there exists an alternating Hamilton path on  $R \cup B$  which begins with  $r$  and ends with another point of  $R$ , and which has at most  $|R| - 2 = |B| - 1$  crossings. Furthermore, if a point, say  $s$ , of  $R \setminus \{r\}$  is a vertex of  $\text{conv}((R \cup B) \setminus \{r\})$ , then there exists an alternating Hamilton  $r - s$  path which has at most  $|R| - 2 = |B| - 1$  crossings. Note that  $s$  is not necessary to be a vertex of  $\text{conv}(R \cup B)$ .*

We now prove our main theorem using the above lemma.

*Proof of Theorem 1.* Suppose that one vertex of  $\text{conv}(X \cup Y)$  is contained in  $X$  and another vertex of  $\text{conv}(X \cup Y)$  is contained in  $Y$ . Then we can choose two consecutive vertices  $x$  and  $y$  of  $\text{conv}(X \cup Y)$  so that  $x \in X$  and  $y \in Y$ . By Lemma 2, there exists an alternating Hamilton  $x - y$  path  $P(x, y)$  which has at most  $|X| - 1$  crossings. Therefore  $P(x, y) + xy$  is the desired alternating Hamilton cycle.

We next assume that all the vertices of  $\text{conv}(X \cup Y)$  are contained in  $X$ . Let  $x \in X$  be a vertex of  $\text{conv}(X \cup Y)$ . Then  $X \setminus \{x\}$ ,  $Y$  and  $x$  satisfy the assumption of Lemma 1, and so there exists a partition  $D_1 \cup D_2$  of  $(X \setminus \{x\}) \cup Y$  which satisfies the conditions in Lemma 1. In particular, for every  $i \in \{1, 2\}$ , we can take a vertex  $d_i$  of  $\text{conv}(D_i)$  which is contained in  $Y$  and adjacent to  $x$ . Then by applying Lemma 2 to  $D_2 \cup \{x\}$ , we can draw an alternating Hamilton path  $P(x, d_2)$  on  $D_2 \cup \{x\}$  which has at most  $|D_2 \cap Y| - 1$  crossings. Similarly by applying Lemma 2 to  $D_1 \cup \{d_2\}$ , in which  $d_1$  is a vertex of  $\text{conv}((D_1 \cup \{d_2\}) \setminus \{d_2\})$ , we can obtain an alternating Hamilton path  $P(d_2, d_1)$  on  $D_1 \cup \{d_2\}$  that has at most  $|D_1 \cap Y| - 1$  crossings. Therefore we can obtain the desired alternating Hamilton cycle  $P(x, d_2) + P(d_2, d_1) + d_1x$ , which has at most  $|D_2 \cap Y| - 1 + |D_1 \cap Y| - 1 + 1 = |Y| - 1$  crossings.  $\square$

*Proof of Lemma 2* We prove the two statements of Lemma 2 at the same time

by induction on  $|R \cup B|$ . The cases in which  $|R \cup B| \leq 4$  are trivial. So we may assume  $|R \cup B| \geq 5$ . For convenience, we color all the points of  $R$  with *red* and those of  $B$  with *blue*.

We now prove Statement (i). First suppose that all the vertices of  $\text{conv}(R \cup B)$  are red. Then by Lemma 1, there exists a partition  $D_1 \cup D_2$  of  $(R \setminus \{r\}) \cup B$  with respect to  $r$ . We denote by  $d_1 \in D_1 \cap B$  and  $d_2 \in D_2 \cap B$  the two blue points satisfying the condition (iii) of Lemma 1. By applying the inductive hypothesis of (ii) to  $D_1 \cup \{d_2\}$ , we can draw an alternating Hamilton path  $P(d_2, d_1)$  on  $D_1 \cup \{d_2\}$  which has at most  $|D_1 \cap R| - 1$  crossings. Similarly by the inductive hypothesis of (i), there exists an alternating Hamilton path  $P(r, d_2)$  on  $D_2 \cup \{r\}$  which has at most  $|D_2 \cap R|$  crossings. Thus we can obtain an alternating Hamilton  $r - d_1$  path  $P(r, d_2) + P(d_2, d_1)$ , which has at most  $|D_1 \cap R| - 1 + |D_2 \cap R| = |R| - 2$  crossings.

We next assume that at least one vertex of  $\text{conv}(R \cup B)$  is blue, and let  $b$  be such a blue vertex. It suffices to show that we can draw an alternating Hamilton  $r - b$  path which has at most  $|R| - 1$  crossings.

Suppose that there exists a blue vertex, say  $b_1$ , of  $\text{conv}((R \cup B) \setminus \{r\})$  which is different from  $b$  and is visible from  $r$  (i.e., the straight-line segment  $rb_1$  intersects  $\text{conv}((R \cup B) \setminus \{r\})$  only at  $b_1$ ). Then by the inductive hypothesis, we can draw an alternating Hamilton path  $P(b_1, b)$  on  $(R \cup B) \setminus \{r\}$  which has at most  $|B| - 2$  crossings. Therefore we obtain the desired alternating Hamilton  $r - b$  path  $P(b_1, b) + rb_1$  on  $R \cup B$ . Similarly, if there exists a red vertex of  $\text{conv}((R \cup B) \setminus \{b\})$  which is not  $r$  and is visible from  $b$ , then we can get the desired alternating Hamilton  $r - b$  path. Therefore we may assume that all the vertices of  $\text{conv}((R \cup B) \setminus \{r\})$  visible from  $r$  are red except  $b$ , and that a similar situation occurs for  $b$  and  $(R \cup B) \setminus \{r\}$ .

We consider two cases.

*Case 1.*  $r$  and  $b$  are consecutive vertices of  $\text{conv}(R \cup B)$ .

Since  $r$  and  $b$  are adjacent, we can find a red vertex  $r_1$  and a blue vertex  $b_1$  of  $\text{conv}((R \cup B) \setminus \{r, b\})$  which are visible from both  $b$  and  $r$ . Then by induction, we can draw an alternating Hamilton path  $P(b_1, r_1)$  on  $(R \cup B) \setminus \{r, b\}$  which has at most  $|R| - 2$  crossings. Consequently we can get an alternating Hamilton  $r - b$  path  $P(b_1, r_1) + b_1r + r_1b$ , which has at most  $|R| - 1$  crossings since  $b_1r$  and  $r_1b$  may intersect.

*Case 2.*  $r$  and  $b$  are not consecutive vertices of  $\text{conv}(R \cup B)$ .

In this case, there exists a partition  $D_1 \cup D_2$  of  $(R \setminus \{r\}) \cup B$  with respect to  $r$  which satisfies the conditions of Lemma 1. For every  $i \in \{1, 2\}$ , let  $d_i$  denote the blue vertex of  $D_i$  satisfying (iii) of Lemma 1.

If  $b \in D_1$ , then we can obtain the desired alternating Hamilton  $r - b$  path from two alternating Hamilton paths  $P(d_2, b)$  on  $D_1 \cup \{d_2\}$  and  $P(r, d_2)$  on  $D_2 \cup \{r\}$ . Hence we may assume that  $b \in D_2$ .

If  $b \neq d_2$ , then we can obtain the desired alternating Hamilton  $r - b$  path from two alternating Hamilton paths  $P(d_2, d_1)$  on  $D_1 \cup \{d_2\}$  and  $P(d_2, b)$  on  $D_2$  by adding  $rd_1$ . Suppose next  $b = d_2 \in D_2$ . Then both  $d_1$  and  $r$  are vertices of  $\text{conv}((D_2 \setminus \{b\}) \cup \{d_1, r\})$ , and so there exists an alternating Hamilton path  $P(r, d_1)$  on  $(D_2 \setminus \{b\}) \cup \{d_1, r\}$  which has at most  $|D_2 \cap B| - 1$  crossings. Similarly

there exists an alternating Hamilton path  $P(b, d_1)$  on  $D_1 \cup \{b\}$ , which has at most  $|D_1 \cap B| - 1$  crossings. Therefore we obtain the desired alternating Hamilton  $r - b$  path  $P(r, d_1) + P(b, d_1)$ , which has at most  $|D_2 \cap B| - 1 + |D_1 \cap B| - 1 + 1 = |B| - 1$  crossings since  $P(r, d_1)$  and  $P(b, d_1)$  intersect at most once. Consequently the proof of (i) is complete.

Next we prove Statement (ii). Note that if the red point  $s$  exists, then it suffices to show that there exists an alternating Hamilton  $r - s$  path. Assume first that a blue point, say  $b'$ , is a vertex of  $\text{conv}(R \cup B)$  and is adjacent to  $r$ . Then by the inductive hypothesis on  $(R \cup B) \setminus \{r\}$ , we can draw an alternating Hamilton path on  $(R \cup B) \setminus \{r\}$  which begins with  $b'$  and ends with the red point  $s$  if any, or a red point of  $R \setminus \{r\}$  and which has at most  $|B| - 1$  crossings. By adding a straight-line segment  $rb'$  to this path, we obtain the desired alternating Hamilton  $r - s$  path or the desired alternating Hamilton path connecting  $r$  and a red point. Hence we may assume that the two vertices of  $\text{conv}(R \cup B)$  adjacent to  $r$  are red points.

By Lemma 1, there exists a partition  $(D_1, D_2)$  of  $(R \cup B) \setminus \{r\}$ . Let  $d_1$  and  $d_2$  are blue points of  $D_1$  and  $D_2$ , respectively, satisfying the condition (iii) of Lemma 1. Note that an equality  $|D_2 \cap R| = |D_2 \cap B|$  also holds, and thus without loss of generality, we may assume that  $s \in D_1$  if any. Then we can draw an alternating Hamilton path  $P_1$  on  $D_1$  which begins with  $d_1$  and ends with  $s$  if any, or a red point of  $(D_1 \cap R) \setminus \{r\}$  and which has at most  $|D_1 \cap B| - 1$  crossings. Similarly, there exists an alternating Hamilton path  $P(d_1, d_2)$  on  $D_2 \cup \{d_1\}$  which has at most  $|D_2 \cap B| - 1$  crossings. Therefore we can get the desired alternating Hamilton path  $P_1 + P(d_1, d_2) + d_2r$ , which has at most  $|D_1 \cap B| - 1 + |D_2 \cap B| - 1 + 1 = |B| - 1$  crossings.

Consequently the proof of Theorem 1 is complete.  $\square$

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